

The Local-to-global Theorem for Cocycle Expansion

Siqi Liu

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1 The local-to-global theorem for cocycle expanders

Constructing coboundary expanders is notoriously difficult, primarily because it is hard to design bounded-degree complexes that satisfy the condition $B^i = Z^i$. However, for many applications this requirement can be relaxed. Complexes that meet these weaker expansion conditions are known as cocycle expanders.

Definition 1.1. A simplicial complex X is a $(i + 1)$ -dimensional η -cocycle expander if

$$\min_{f \in C^i \setminus Z^i} \frac{|\delta_i f|}{|f - Z^i|} \geq \eta.$$

Analogous to how trickle-down theorem tells us that expansion of the links implies expansion of the global graph, it turns out that coboundary expansion on the links together with local-spectral expansion on the entire complex implies cocycle expansion. This is captured by the following theorem.

Theorem 1.2 ([KKL14, EK24, DD24]). *For any constants $\beta > 0$ and $\gamma > 0$ sufficiently small as a function of β , if X is a 3-dimensional simplicial complex with bounded vertex degree D ¹ such that*

1. X is a 3-dimensional γ -local spectral expander and
2. for all $v \in X(0)$, the link X_v is a 2-dimensional β -coboundary expander,

then X is a 2-dimensional η -cocycle expander for $\eta = \min\left(\frac{1}{D}, \frac{\beta}{17}\right)$.

This result generalize to higher dimensional cocycle-expansion, but in this note we focus on 2-dimensional setting. The proof we present here combines ideas from [EK24, DD24].

1.1 The local correction algorithm

The proof of this theorem uses the local correction argument. The gist is that given any function $f \in C^1$, we iteratively make local corrections to f so that at each step the coboundary size $|\delta_1 f|$ decreases while the correction function has bounded weight. This procedure ends when we get a locally minimal function, defined as follows.

Definition 1.3. A function $f \in C^1$ is locally minimal if for every vertex $v \in X(0)$ and any localized function $h_v : X(1) \rightarrow \mathbb{F}_2$ such that $h_v(e) = 0$, $\forall e \not\ni v$ (in other words h_v can be nonzero only over v 's adjacent edges), f satisfies that

$$|\delta_1 f| \leq |\delta_1(f + h_v)|.$$

The following observation of locally minimal function will be key in our proof.

¹meaning every v is in at most D 3-faces.

Observation 1.4. For any function $f \in C^1(X)$ and any vertex $v \in X(0)$, we define the induced function $f_v \in C^1(X_v)$ to be

$$f_v(uw) = f(uw), \quad \forall uw \in X_v(1)$$

and the projection function $g \in C^0(X_v)$ to be

$$g_v(u) = f(uv), \quad \forall u \in X_v(0).$$

If f is locally minimal, then the local functions satisfy that

$$|f_v - B_v^1| = |f_v - \delta_0 g_v|,$$

where B_v^1 denote the 1-dimensional coboundary space of X_v .

In other words, $\delta_0 g_v$ is a closest function to f_v in B_v^1 .

Now we formally define the algorithm that locally correct any function in C^1 to a locally minimal function.

The local correction algorithm Let $f \in C^1$ be the input. Set $f_0 = f$

At iteration i , if there exists $v_i \in X(0)$ and $h_i : X(1) \rightarrow \mathbb{F}_2$, which is only nonzero over adjacent edges of v_i , such that $|\delta_1(f_i + h_i)| < |\delta_1 f_i|$, update $f_{i+1} = f_i + h_i$. Repeat this until no such v_i, h_i exist.

First note that this process terminates since in every iteration the number of failed triangle test decreases. Let T be the number of iterations on input f , so the final output is f_T .

Two cases If the algorithm outputs a function $f_T \in Z^1$, then by the bound on the distance $\|f - f_T\|$ above we deduce that

$$\frac{|\delta_1 f|}{|f - Z^1|} \geq \frac{\sum_{i=1}^T |\delta_1 f_i| - |\delta_1 f_{i-1}|}{\sum_{i=1}^T \|h_i\|} \geq \min_i \frac{\min_{ab \in X_{v_i}(1)} \pi_2(v_i ab)}{\pi_1(\{e \in X(1) \mid e \ni v_i\})} \geq \frac{1}{D}.$$

Here the last equality holds for all bounded-degree simplicial complex with uniform distribution on their maximal faces. In fact, a more careful argument can remove this assumption and yield $\Theta(\beta)$ using coboundary expansion of X_v 's. We omit it here and you will prove it in the exercise.

If $f_T \notin Z^1$, then we shall prove that in this case:

Claim 1.5. For any simplicial complex X that satisfies the conditions in Theorem 1.2, any locally minimal function $f \in C^1 \setminus Z^1$ has $|\delta_1 f| > \frac{\beta}{17}$.

Using this lower bound for the failure probability of locally minimal functions, we can derive that in this case

$$\frac{|\delta_1 f|}{|f - Z^1|} > \frac{\beta}{17}.$$

From the two case analysis we can conclude that X is a 2-dimensional η -cocycle expander, where $\eta = \min\left(\frac{1}{D}, \frac{\beta}{17}\right)$.

1.2 Proof of Claim 1.5

Given a locally minimal $f \in C^1 \setminus Z^1$, denote by $T^* = \{t \in X(2) \mid \delta_1 f(t) \neq 0\}$ the set of failed triangle tests of f . We shall prove that assume for the sake contradiction, $|\delta_1 f| = \|\vec{1}_{T^*}\|^2 \leq \frac{\beta}{17}$ then there is an operator $P : \ell^2(X(2)) \rightarrow \ell^2(X(2))$ with $\|P\|_{op} \leq \gamma$ so that

$$\left(\frac{1}{12} - O\left(\frac{\gamma^2}{\beta^2}\right) - O(\gamma^2)\right) \cdot \|\vec{1}_{T^*}\|^2 \leq \langle \vec{1}_{T^*}, P \vec{1}_{T^*} \rangle. \quad (1)$$

Then by the assumption that γ is small we can derive the contradiction that

$$\left(\frac{1}{12} - O\left(\frac{\gamma^2}{\beta^2}\right) - O(\gamma^2)\right) \cdot \|\vec{1}_{T^*}\|^2 \leq \langle \vec{1}_{T^*}, P\vec{1}_{T^*} \rangle \leq \gamma \cdot \|\vec{1}_{T^*}\|^2.$$

To prove (1), we set $P = M_2^+ - U_1 D_2$, where M_2^+ is the non-lazy up-down walk over triangles and $U_1 D_2$ is the standard down-up walk on triangles. First note that by the γ -local spectral expansion of X , the norm of their difference is indeed $\|M_2^+ - U_1 D_2\|_{op} \leq \gamma$ (Lemma A.3).

Next, we notice that $\vec{1}_{T^*} = \delta_1 f$. So for any $\sigma \in X(3)$ we have $\delta_2 \vec{1}_{T^*}(\sigma) = 0$. In other words, if $t \in T^*$, then for any $\sigma \in X(3)$ that contains t , at least one of the other 3 triangles in σ must also be in T^* . Therefore we can derive

$$\langle \vec{1}_{T^*}, M_2^+ \vec{1}_{T^*} \rangle \geq \frac{1}{3} \cdot \|\vec{1}_{T^*}\|^2.$$

We next prove that

$$\langle \vec{1}_{T^*}, U_1 D_2 \vec{1}_{T^*} \rangle \leq \left(\frac{1}{4} + O\left(\frac{\gamma^2}{\beta^2}\right) + O(\gamma^2)\right) \cdot \|\vec{1}_{T^*}\|^2.$$

We rewrite the inner product as follows.

$$\begin{aligned} \langle \vec{1}_{T^*}, U D \vec{1}_{T^*} \rangle &= \mathbb{E}_{ab \sim \pi(1)} \left[\Pr_{c, d \sim \pi_{ab}(0)} [abc, abd \in T^*] \right] \\ &= \mathbb{E}_{a \sim \pi(0), b \sim \pi_a(0)} [\mathbb{E}_{t \supset ab} [\delta_1 f(t)^2]]. \end{aligned}$$

We call an edge ab heavy if the local failure probability $\mathbb{E}_{t \supset ab} [\delta_1 f(t)]$ is large ($\geq \frac{1}{4}$). Similarly, we call a vertex a heavy if the local failure probability $\mathbb{E}_{t \ni a} [\delta_1 f(t)]$ is large ($\geq \frac{1}{16}$).

To bound the inner product, we need to show that the edge failure probability concentrates well around the expectation ε , so there aren't many heavy edges. An edge ab could be heavy for one of the two reasons: 1. either a or b is a heavy vertex; 2. or despite a and b both not heavy, the edge ab does not "sample" the failure probability at the vertices well.

Specifically, denote the first bad event to be

$$E_1 = \left\{ ab \in X(1) \mid \Pr_{t \ni a} [\delta_1 f(t) \neq 0] \geq \frac{1}{16} \text{ or } \Pr_{t \ni b} [\delta_1 f(t) \neq 0] \geq \frac{1}{16} \right\},$$

and the second bad event to be

$$E_2 = \left\{ ab \in X(1) \mid ab \notin E_1 \text{ but } \Pr_{t \supset ab} [\delta_1 f(t) \neq 0] \geq \frac{1}{4} \right\}.$$

If ab is in neither E_1 or E_2 , we have that $\Pr_{t \supset ab} [\delta_1 f(t) \neq 0] < \frac{1}{4}$.

So now we bound the inner product by

$$\begin{aligned} \langle \vec{1}_{T^*}, U D \vec{1}_{T^*} \rangle &\leq \Pr_{ab \sim \pi(1)} [E_1] + \Pr_{ab \sim \pi(1)} [E_2] \\ &\quad + \mathbb{E}_{ab \sim \pi(1)} \left[\mathbf{1}[\neg E_1 \wedge \neg E_2] \cdot \frac{1}{4} \mathbb{E}_{t \supset ab} [\delta_1 f(t)] \right] \\ &\leq \Pr_{ab \sim \pi(1)} [E_1] + \Pr_{ab \sim \pi(1)} [E_2] + \frac{1}{4} \|\vec{1}_{T^*}\|^2 \end{aligned}$$

So we are left to bound the probability of E_1 and E_2 for locally minimal f .

Claim 1.6. For an X that satisfies the conditions in Theorem 1.2, any locally minimal $f \in C^1 \setminus Z^1$ with $\|\vec{1}_{T^*}\|^2 := \varepsilon \leq \frac{\beta}{17}$ satisfies that

$$\Pr_{ab \sim \pi(1)} [ab \in E_1] \leq \varepsilon \cdot O\left(\frac{\gamma^2}{\beta^2}\right).$$

Proof. First define the set of heavy vertices to be

$$V^* = \left\{ v \in X(0) \mid \Pr_{t \ni v} [\delta_1 f(t) \neq 0] \geq \frac{1}{16} \right\}.$$

We define another related set

$$R = \left\{ v \in X(0) \mid \Pr_{t \in X_v(2)} [\delta_1 f(t) \neq 0] \geq \frac{\beta}{16} \right\}.$$

We shall first show that $R \supseteq V^*$ and then bound the weight of R .

For every $v \in V^*$ define the induced local functions $g_v : X_v(0) \rightarrow \mathbb{F}_2$ as

$$g_v(a) = f(\{a, v\}).$$

By construction in the link of v

$$\delta_0 g_v(\{a, b\}) \neq f(\{a, b\}) \Rightarrow f(\{v, a\}) + f(\{v, b\}) + f(\{a, b\}) \neq 0 \Rightarrow \delta_1 f(\{a, b, v\}) \neq 0.$$

Furthermore by local minimality of f , $|f_v - B_v^1| = |f_v - \delta_0 g_v| = \Pr_{t \ni v} [\delta_1 f(t) \neq 0]$. Since X_v is a 2-dimensional β -coboundary expander,

$$\frac{|\delta_1 f_v|}{|f_v - B_v^1|} \geq \beta \Rightarrow |\delta_1 f_v| \geq \beta \cdot \Pr_{t \ni v} [\delta_1 f(t) \neq 0].$$

So for every $v \in V^*$, we have that

$$\Pr_{t \in X_v(2)} [\delta_1 f(t) \neq 0] \geq \beta \cdot \Pr_{t \ni v} [\delta_1 f(t) \neq 0] \geq \frac{\beta}{16}.$$

Therefore $v \in R$ as well.

We next bound the probability that a random vertex is in R . Consider the bipartite graph G between $X(0)$ and $X(2)$ where we connect v and t if and only if $\{v\} \cup t \in X(3)$. Use σ_2 to denote the second largest singular value of the normalized biadjacency matrix of G . Then by the bipartite expander sampling lemma (Lemma A.1) we have that

$$\Pr_{v \sim \pi(0)} \left[\Pr_{t \in X_v(2)} [\delta_1 f(t) \neq 0] \geq \frac{\beta}{16} \right] \leq \varepsilon \cdot \frac{\sigma_2^2}{\left(\frac{\beta}{16} - \varepsilon\right)^2}$$

We thereby deduce that

$$\Pr_{v \sim \pi(0)} [v \in V^*] \leq \Pr_{v \sim \pi(0)} [v \in R] \leq \varepsilon \cdot \frac{\sigma_2^2}{(\beta/16 - \varepsilon)^2}.$$

By the spectral bound for the swap walks covered in Max's lectures, know that $\sigma_2^2 \leq 3\gamma^2$ (Lemma A.2). Thus we conclude by observing

$$\Pr_{ab \sim \pi(1)} [ab \in E_1] \leq 2 \Pr_{v \sim \pi(0)} [v \in R] \leq \varepsilon \cdot O\left(\frac{\gamma^2}{\beta^2}\right).$$

□

The other claim needed in step 2 is the following.

Claim 1.7. *For an X that satisfies the conditions in Theorem 1.2, any locally minimal $f \in C^1 \setminus Z^1$ with $\|\vec{1}_{T^*}\|^2 = \varepsilon$ satisfies that*

$$\Pr_{ab \sim \pi(1)} [E_2] \leq \varepsilon \cdot O(\gamma^2).$$

Proof. For any edge $ab \in E_2$ we have that by the local minimality of f , the local failure probability at ab is $\Pr_{t \supset ab} [\delta_1 f(t)] \in [\frac{1}{4}, \frac{1}{2}]$. So in the link of ab , the set of triangles $T^* \cap \{t \in X(2) \mid t \supset ab\}$ induces a large cut C_{ab} in the link of ab . By the expander mixing lemma, the C_{ab} contains at least $\frac{1}{4} \cdot \frac{3}{4} - \gamma$ -fraction of edges. For any $cd \in C_{ab}$, we know that exactly one of abc and abd are in T^* , so exactly one of acd and bcd must be in T^* as well. Wlog assume that for at least half of the $cd \in C_{ab}$, we have $acd \in T^*$, i.e.

$$\Pr_{cd \in X_{ab}(1)} [acd \in T^*] \geq \frac{3}{32} - \frac{\gamma}{2},$$

while

$$\Pr_{cd \in X_a(1)} [acd \in T^*] < \frac{1}{16}.$$

For such a fixed vertex a , to bound the probability $\Pr_{b \in X_a(0)} [\Pr_{cd \in X_{ab}(1)} [acd \in T^*] \geq \frac{3}{32} - \frac{\gamma}{2}]$, we consider the bipartite graph G_a between $X_a(0)$ and $X_a(1)$ where we connect $b \in X_a(0)$ and $cd \in X_a(1)$ if and only if $bcd \in X_a(2)$. Use σ_a to denote the second largest singular value of the normalized biadjacency matrix of G_a . Then by the bipartite graph expander sampling lemma we get that

$$\Pr_{b \in X_a(0)} \left[\Pr_{cd \in X_{ab}(1)} [acd \in T^*] \geq \frac{3}{32} - \frac{\gamma}{2} \right] \leq \Pr_{cd \in X_a(1)} [acd \in T^*] \cdot \frac{\sigma_a^2}{\left(\frac{1}{32} - \frac{\gamma}{2}\right)^2}.$$

Averaging over all non-heavy vertices a we obtain

$$\begin{aligned} \Pr_{ab \sim \pi(1)} [E_2] &\leq \mathbb{E}_{a \sim \pi(0)} \left[\Pr_{b \in X_a(1)} [ab \text{ is heavy}] \cdot \vec{1}[a, b \text{ not heavy}] \cdot \vec{1} \left[\Pr_{cd \in X_{ab}(1)} [acd \in T^*] \geq \frac{3}{32} - \frac{\gamma}{2} \right] \right] \\ &\quad + \mathbb{E}_{a \sim \pi(0)} \left[\Pr_{b \in X_a(1)} [ab \text{ is heavy}] \cdot \vec{1}[a, b \text{ not heavy}] \cdot \vec{1} \left[\Pr_{cd \in X_{ab}(1)} [bcd \in T^*] \geq \frac{3}{32} - \frac{\gamma}{2} \right] \right] \\ &\leq 2 \mathbb{E}_{a \sim \pi(0)} \left[\Pr_{b \in X_a(1)} \left[\vec{1}[a \text{ not heavy}] \cdot \Pr_{cd \in X_{ab}(1)} [acd \in T^*] \geq \frac{3}{32} - \frac{\gamma}{2} \right] \right] \\ &\leq 2 \mathbb{E}_{a \sim \pi(0)} \left[\Pr_{cd \in X_a(1)} [acd \in T^*] \right] \cdot \frac{\max_{a \in X(0)} \sigma_a^2}{\left(\frac{1}{32} - \frac{\gamma}{2}\right)^2} \\ &= \varepsilon \cdot O(\gamma^2) \end{aligned}$$

The last inequality follows from the spectral expansion of swap walks in γ -local spectral expanders (Lemma A.2). $\square$

A Auxiliary lemmas

In the proof we used the following results. We omit their proofs because they are either standard or have been covered in an earlier talk.

Lemma A.1. *Let $G = (L, R, E)$ be a biregular bipartite graph whose normalized biadjacency operator has second singular value at most σ . For every $f : L \rightarrow \mathbb{R}$, let*

$$\mu = \mathbb{E}_{v \sim L} [f(v)].$$

Then, for every $\varepsilon > 0$,

$$\Pr_{u \sim R} [\mathbb{E}_{v \sim N(u)}[f(v)] > \mu + \varepsilon] \leq \frac{\sigma^2 \text{Var}_{v \sim L}[f(v)]}{\varepsilon^2}.$$

Lemma A.2 ([DD19]). *Let X be a d -dimensional (two-sided) λ -local spectral expander. Let $i, j \geq 0$ satisfy $i + j + 1 \leq d$, and let*

$$S_{i,j} : L^2(X(i), \pi_i) \longrightarrow L^2(X(j), \pi_j)$$

be the normalized (i, j) -swap walk operator. Then

$$\sigma_2(S_{i,j}) \leq (i+1)(j+1)\lambda.$$

Lemma A.3 ([DDFH24]). *Let X be a d -dimensional (two-sided) λ -local spectral expander. Then, for every $0 \leq i \leq d-1$,*

$$\|M_i^+ - U_{i-1}D_i\|_{op} \leq \lambda,$$

where M_i^+ is the non-lazy up-down walk on $X(i)$, and the operator norm is with respect to $L^2(X(i), \pi_i)$.

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